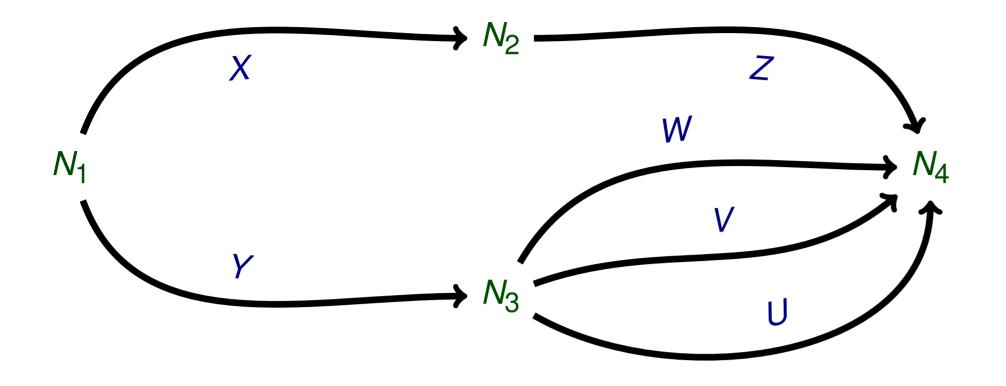
Dynamic Programming and Subtree Perfectness for Deterministic Discrete-Time Systems with Uncertain Rewards

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Background

Backward Induction

Backward Induction Method

apply choice function recursively to gambles from last stage

- first, apply opt on $\{W, V, U\}$, say opt $(\{W, V, U\}) = \{V, U\}$
- next, only need opt on $\{X + Z, Y + W, Y + V, Y + U\}$

Problem

- will not always coincide with the standard normal form solution
- e.g. Γ-maximin and interval dominance do not [1]
- partial answer in [1]: sufficient condition for backward induction under assumption of utility

Normal Form Solutions

aim: find optimal paths from N_1 to N_4 set of paths judged optimal is a normal form solution

Rewards

■ binary operator +: for combining rewards. left identity element: 0 + r = rleft inverse: (-r) + r = 0no additional structure assumed (no utility, no full ranking)

Gambles and Choice Functions

each path corresponds to a gamble (an uncertain reward)

 $N_1 \rightarrow N_2 \rightarrow N_4 = X + Z$

find optimal paths by finding optimal gambles choice function opt maps sets of gambles to (optimal) subsets

Standard Normal Form Solution

apply choice function to the set of all the problem's gambles choose one of the paths judged optimal by this process

> inefficient for large problems! but we can do more cleverly...

References

[1] G. De Cooman and M.C.M. Troffaes.

Dynamic programming for deterministic discrete-time systems with uncertain

- our contribution:

necessary and sufficient conditions for backward induction necessary and sufficient conditions for subtree perfectness no utility assumed

Necessary and Sufficient Conditions for Backward Induction

Insensitivity to Omission of Non-Optimal Elements

 $opt(\mathcal{X}) \subseteq \mathcal{Y} \subseteq \mathcal{X} \Rightarrow opt(\mathcal{Y}) = opt(\mathcal{X}).$

Preservation of Non-Optimality Under Addition of Elements

 $\mathcal{Y} \subseteq \mathcal{X} \Rightarrow \mathsf{opt}(\mathcal{Y}) \supseteq \mathsf{opt}(\mathcal{X}) \cap \mathcal{Y}.$

Backward Addition Property

 $opt(X + Y) \subseteq X + opt(Y).$

Subtree Perfectness

Example of Failure of Subtree Perfectness

- say opt($\{W, V, U\} = \{V, U\}$, so W is deleted
- but global solution is $\{X + Z, Y + U\}$, so also V is deleted
- removal of V not caused by the local choice at N_3 but by global choice when X + Z was also considered at N_1
- choice at N_3 is not completely determined by the options available at N_3
- this is a failure of subtree perfectness: choice in subproblem is influenced by problem into which it is embedded [3]

gain.

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International Journal of Game Theory, 4(1):25–55, Mar 1975.

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Necessary and Sufficient Conditions for Subtree Perfectness

Intersection Property

 $\mathsf{opt}(\mathcal{Y}) = \mathsf{opt}(\mathcal{X}) \cap \mathcal{Y}.$

Addition Property

$$opt(X + Y) = X + opt(Y)$$

Final Remarks

Insensitivity to Omission implied by Intersection

Intersection is equivalent to Total Preorder

indeterminate choice functions can never be subtree perfect similar results well known for other decision problems [2, 4]



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